



TRUTH TABLE

X ₁	X ₀	Y ₁	Y ₀	K	2 ¹	2 ²	2 ³	2 ⁴	2 ⁵	2 ⁶
0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	1	0	0	0	0	0	0
0	0	0	1	0	0	0	0	0	0	0
0	0	0	1	1	0	0	0	0	0	0
0	1	0	0	0	0	0	0	0	0	0
0	1	0	0	1	0	0	0	0	0	0
0	1	0	1	0	0	0	0	0	0	0
0	1	0	1	1	0	0	0	0	0	0
1	0	0	0	0	0	0	0	0	0	0
1	0	0	0	1	0	0	0	0	0	0
1	0	0	1	0	0	0	0	0	0	0
1	0	0	1	1	0	0	0	0	0	0
1	1	0	0	0	0	0	0	0	0	0
1	1	0	0	1	0	0	0	0	0	0
1	1	0	1	0	0	0	0	0	0	0
1	1	0	1	1	0	0	0	0	0	0

V_{CC} = Pin 24
 V_{EE} = Pin 12

P_D = 760 mW typ/pkg (No Load)
 t_{pd} = 50 ns typ (8 x 8 bit product)

4 x 2 Multiplier

The MC10183 is a 4 x 2 bit multiplier that can multiply 2's complement numbers producing a 2's complement product without correction. The device can be used as a 4 x 2 bit multiplier cell to build larger iterative arrays.

The part performs the function defined as $F = XY + K$, where K is an input field used to add partial products in an array or to add a constant to the least significant part of the array product. The algorithm used is a modified Booth's algorithm or multiplier coding technique.

The device consists of a shift network and an adder/subtractor in which 0, 1 times X, or 2 times X is either added or subtracted to input constant K. The Y inputs control multiplication as shown in the Truth Table. The most significant digit in a word carries a negative weight allowing 2's complement numbers of various lengths to be multiplied. An M-bit by N-bit multiplication produces an M + N bit product.

The \bar{P} polarity input allows multiplication in either positive logic ($\bar{P} = \text{high}$) or negative logic ($\bar{P} = \text{low}$) representation. Also, mode control M inverts \bar{C}_N when high and passes \bar{C}_N directly when left low.

MC10183

ARITHMETIC FUNCTIONS